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BOUNDING COCHORDAL COVER NUMBER OF GRAPHS VIA VERTEX STRETCHING

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ABSTRACT. In this paper, it is shown that when a special vertex stretching is applied to a graph, the cochordal cover number of the graph increases at most two. As a consequence, it is shown that the induced matching number and cochordal cover number of a special vertex stretching of a graph G are equal provided G is well-covered bipartite or weakly chordal graph.

Keywords: Castelnuovo-Mumford regularity, induced matching number, cochordal cover number.

MSC(2010): Primary: 13F55; Secondary: 05C70, 05E45.

1. Introduction

Let G be a simple graph with vertex set $V(G) = \{1, \dots, n\}$ and edge set $E(G)$. A *matching* in a graph G is a set of edges M such that no two edges share a common end. If M is an induced subgraph, the matching is called an *induced matching*. The maximum size of an induced matching in G is called the *induced matching number* of G and denoted by $\text{indmatch}(G)$. The problem of finding a maximum induced matching is known to be NP-hard.

Let G be a graph and x be a vertex in G . A *vertex stretching* with respect to x is defined to be the transformation consisting of the following steps:

- (1) partition the open neighborhood $N(x)$ of vertex x into two disjoint subsets Y and Z in an arbitrary way;
- (2) delete vertex x from the graph together with incident edges;
- (3) add a chordless path on k vertices $(a, x_1, \dots, x_{k-2}, b)$ to the remaining graph;
- (4) connect a to each vertex in Y , and connect b to each vertex in Z .

Notice that k is a parameter associated with the transformation. We shall denote the vertex stretching with parameter k by \mathcal{P}^k . The graph produced by this operation will be denoted $\mathcal{P}^k(G) = \mathcal{P}^k(G, x)$.

Vertex stretching have been applied to different problems for several goals (see e.g. [1, 4, 8, 12] and [2]).

Although the problem of finding a maximum induced matching is known to be NP-hard, in [11] Lozin showed that \mathcal{P}^4 increases the size of a maximum induced matching by one. As a consequence, Lozin showed that the induced matching problem remains NP-hard in some special classes of bipartite graphs such as bipartite graphs with maximum degree 3 or C_4 -free bipartite graphs.

Let x and Y be non adjacent vertices of a graph G . *Identify x and y* means to replace these vertices by a single vertex incident to all the edges which were incident in G to either x or y . We denote the resulting graph by $G/\{x, y\}$.

Let $R = \mathbb{K}[x_1, \dots, x_n]$ be the polynomial ring over a field \mathbb{K} . The *edge ideal* of G in R is defined by $\mathcal{I}(G) = (x_i x_j \mid \{i, j\} \in E(G))$. In the recent years, several authors have studied the Castelnuovo–Mumford regularity of $R/\mathcal{I}(G)$ denoted by $\text{reg}(R/\mathcal{I}(G))$. It is known that $\text{indmatch}(G) \leq \text{reg}(R/\mathcal{I}(G))$, cf. [7, Lemma 2.2]. Recently, Biyikoglu and Civan [3] get motivated to show that \mathcal{P}^4 increases the size of the Castelnuovo–Mumford regularity by one.

In [13] Woodroffe introduced the notion of cochordal cover number of a graph to study the Castelnuovo–Mumford regularity of a graph. A graph G is called *chordal* if every induced cycle in G has length 3, and is *cochordal* if the complement graph G^c is chordal. The *cochordal cover number* of G is the minimum number of cochordal subgraphs required to cover the edges of G and it is denoted by $\text{cochord}(G)$.

Combining [7, Lemma 2.2] and [13, Lemma 1] lead to the following inequalities:

$$\text{indmatch}(G) \leq \text{reg}(R/\mathcal{I}(G)) \leq \text{cochord}(G).$$

Now it is natural to ask: what can we say about the cochordal cover number of $\mathcal{P}^4(G)$?. In this note, we show that \mathcal{P}^4 increases the cochordal cover number by at most two. Using [11, Corollary 1] this result implies that the computational complexity of the cochordal cover number of arbitrary graphs is equivalent to that of bipartite graphs having sufficiently large girth with maximum degree three.

2. The results

To prove the main result we need the following lemmas.

Lemma 2.1. *Let G be a graph and x, y be two non-adjacent vertices of G . If $N(x) \cap N(y) = \emptyset$ and G is cochordal graph, then $H = G/\{x, y\}$ is cochordal graph.*

Proof. At first note that every vertex of $N_G(x)$ is adjacent to every vertex of $N_G(y)$, since otherwise if $x' \in N_G(x)$ is not adjacent to $y' \in N_G(y)$ in G , then $x - y - x' - y' - x$ is an induced cycle of length four in G^c , which is a contradiction.

Let $x = y = a$ in graph H . Suppose that C is an induced cycle of length k in $(G/\{x, y\})^c$. If a is not a vertex of C , then C is a cycle of G^c and hence $k = 3$. Suppose that a is a vertex of C and

$$C : a = v_1 - v_2 - \cdots - v_k - v_1 = a.$$

Then $v_2, v_k \notin N_H(a)$ and $v_3, v_4, \dots, v_{k-1} \in N_H(a)$. Since $N_H(a) = N_G(x) \cup N_G(y)$, we conclude that $\{v_2, v_k\} \subseteq N_{G^c}(x) \cap N_{G^c}(y)$ and $v_i \in N_G(x) \cup N_G(y)$. Suppose that $v_3 \in N_G(x)$. Since every vertex of $N_G(x)$ is adjacent to every vertex of $N_G(y)$, we conclude that $v_4 \in N_G(x)$. By the same argument we conclude that $\{v_3, v_4, \dots, v_{k-1}\} \subseteq N_G(x)$. Hence the cycle $x - v_2 - v_3 - \cdots - v_{k-1} - v_k - x$ is an induced cycle of G^c . Thus $k = 3$ and H is a chordal graph. \square

Lemma 2.2. *Let G be a cochordal graph. Then*

$$2 \leq \text{cochord}(\mathcal{P}^4(G)) \leq 3.$$

Proof. Let $N_G(x) = X \cup Y$ be a partition of neighborhoods of x and $\mathcal{P}^4(G) = \mathcal{P}^4(G, x)$ is obtained from G , by deleting the vertex x , adding the path $x_1 - a - b - x_2$, joining the vertex x_1 to vertices in X and joining the vertex x_2 to vertices in Y . Define three subgraph G_1, G_2, G_3 of $\mathcal{P}^4(G)$ as follows:

$$V(G_1) = X \cup \{x_1, a, b\}, E(G_1) = \{x_1a, ab\} \cup \{x_1z : z \in X\},$$

$$V(G_2) = Y \cup \{x_2, b\}, E(G_2) = \{x_2b\} \cup \{x_2z : z \in Y\},$$

$$V(G_3) = V(\mathcal{P}^4(G)) \setminus \{x_1, x_2, a, b\}, E(G_3) = E(\mathcal{P}^4(G)) \setminus (E(G_1) \cup E(G_2)).$$

It is not difficult to see that G_i 's are cochordal edge-disjoint subgraphs $\mathcal{P}^4(G)$ and cover the edges of $\mathcal{P}^4(G)$. Therefore $\text{cochord}(\mathcal{P}^4(G)) \leq 3$. On the other hand if $X \neq \emptyset$ (or $Y \neq \emptyset$), we can consider $z \in X$ ($z \in Y$) and by considering four vertices z, x_1, x_2b (or a) we have a induced cycle C_4 in $\mathcal{P}^4(G)^c$. Hence $2 \leq \text{cochord}(\mathcal{P}^4(G))$. \square

Remark 2.3. Note that for any graph G , the cochord number of $\mathcal{P}^k(G) = \mathcal{P}^k(G, x)$ depends on x and partition of $N(x)$; for example $\mathcal{P}^4(C_4)$ is C_7 or $C_4.P_4$ (this is a graph union of C_4 and P_4 where their intersection has just one vertex which is the end vertex of P_4) and we have $\text{cochord}(C_7) = 3$ and $\text{cochord}(C_4.P_4) = 2$ and we have $\text{cochord}(C_7) = 3$ and $\text{cochord}(C_4.P_4) = 2$.

Lemma 2.4. *Let G be a cochordal graph. Then $\text{cochord}(\mathcal{P}^4(G, x)) = 2$, if one of the following holds:*

- i) One of the partitions of $N(x)$ is empty;
- ii) Induced subgraph $\langle N_G(x) \rangle$ is a clique.

Proof. i) By the same notation as in the proof of Lemma 2.2, Suppose that $Y = \emptyset$. Hence two graphs

$$G_1 : x_1 - a - b - x_2, G_2 : \mathcal{P}^4(G) \setminus \{a, b, x_2\}$$

are edge-disjoint cochordal subgraphs of $\mathcal{P}^4(G)$ which cover the edge set of $\mathcal{P}^4(G)$. Hence $\text{cochord}(\mathcal{P}^4(G)) \leq 2$. Therefore $\text{cochord}(\mathcal{P}^4(G)) = 2$

ii) In this case consider the following edge-disjoint subgraphs of $\mathcal{P}^4(G)$:

$$G_1 : x_1 - a - b - x_2, G_2 : \mathcal{P}^4(G) \setminus \{a, b\}.$$

The subgraph G_1 is a cochordal graph, since G_1 is a path of length three. We prove that the subgraph G_2 is a cochordal graph. Suppose that C is an induced cycle of length k in G_2^c . If $x_1, x_2 \notin V(C)$, then C is an induced subgraph of G^c , and hence $k = 3$. Suppose that $x_1, x_2 \in C$ and hence we can show the cycle C as follows:

$$C : x_1 = v_1 - v_2 - \dots - v_{k-1} - (v_k = x_2) - x_1 = v_1.$$

Therefore

$$\{v_3, v_4, \dots, v_{k-1}\} \subseteq N_G(x_1) = X,$$

$$\{v_2, v_3, \dots, v_{k-2}\} \subseteq N_G(x_2) = Y.$$

But $X \cap Y = \emptyset$, implies that $k \leq 4$. Suppose that $k = 4$. Hence v_2 is adjacent to v_3 in G_2^c , which is a contradiction, since $N_G(x)$ is a clique. Therefore $k = 3$. Now consider the case $|V(C) \cap \{x_1, x_2\}| = 1$. Without lose of generality assume that $x_1 \in V(C)$ and $x_2 \notin V(C)$ (the proof of the case $x_1 \notin V(C)$ and $x_2 \in V(C)$ is similar). Again suppose that the cycle C is as the following form:

$$C : x_1 = v_1 - v_2 - \dots - v_k - v_1 = x_1.$$

Hence

$$\{v_3, v_4, \dots, v_{k-1}\} \subseteq N_G(x_1) = X,$$

Since $N_G(x)$ is a clique and $X \subseteq N_G(x)$, we conclude that $k \leq 4$. If $k = 4$, then $v_3 \in X$ and $v_2, v_4 \notin N_G(x)$. Hence the cycle $x - v_2 - v_3 - v_4 - x$ is an induced cycle of length four in G^c , which is a contradiction. Hence $k = 3$ and therefore G_2 is a cochordal graph. \square

Corollary 2.5. For $n \geq 2$, $\text{cochord}(\mathcal{P}^4(K_n)) = 2$.

Lemma 2.6. Let G be a cochordal graph. If G has an induced cycle C_4 , then there exists a vertex x , such that $\text{cochord}(\mathcal{P}^4(G)) = 3$, for some partitions of $N_G(x)$.

Proof. Let x be a vertex of induced cycle C_4 of G and y be a neighborhood of x among the vertices of cycle C_4 . Let $X = N_G(x) \setminus \{y\}$ and $Y = \{y\}$. By this partition of $N_G(x)$, the graph $\mathcal{P}^4(G)$ has an induced cycle C_7 . If $\text{cochord}(\mathcal{P}^4(G)) = 2$ and G_1 and G_2 are cochrdal, edge-disjoint subgraphs, which cover the edges of $\mathcal{P}^4(G)$, then G_1 or G_2 has an induced path P_5 . Hence one of the subgraphs G_1 or G_2 are not cochordal graph, which is a contradiction. Hence $\text{cochord}(\mathcal{P}^4(G)) = 3$. \square

Theorem 2.7. *Let G be a simple graph. Then*

$$\text{cochord}(G) + 1 \leq \text{cochord}(\mathcal{P}^4(G)).$$

Proof. Let $x \in V(G), Y, Z \subseteq N(x)$ such that $Y \cup Z = N(x)$ and $Y \cap Z = \emptyset$.

Consider four new vertices x_1, x_2, a, b . Suppose that

$$\mathcal{P}^4(G) = (G - x) + \{ay : y \in Y\} + \{bz : z \in Z\} + \{x_1a, x_1x_2, x_2b\}.$$

Let $\text{cochord}(\mathcal{P}^4(G)) = \ell$ and G_1, \dots, G_ℓ be edge disjoint subgraphs of $\mathcal{P}^4(G)$ such that $E(\mathcal{P}^4(G)) = \bigcup_1^\ell E(G_i)$ and G_i^c is chordal graph for any $1 \leq i \leq \ell$. Since $E(G_i)$ s cover the edges of $\mathcal{P}^4(G)$, there exists an i_0 such that $x_2b \in E(G_{i_0})$. If $a \notin V(G_{i_0})$, then $N_{G_{i_0}}(a) = \emptyset$ and if $a \in V(G_{i_0})$, since $G_{i_0}^c$ is chordal, then $N_{G_{i_0}}(a) = \{x_1\}$. In any case $N_{G_{i_0}}(a) \subseteq \{x_1\}$. In addition, if the edges $x_1a, x_2b \in E(G_{i_0})$, then $x_1x_2 \in E(G_{i_0})$. Otherwise the edges $x_1x_2, x_1b, ab, x_2a \in E(G_{i_0}^c)$ and therefore $G_{i_0}^c$ has an induced cycle C_4 , which is a contradiction. Thus the path $a - x_1 - x_2 - b$ is an induced subgraph of G_{i_0} . Since $G_{i_0}^c$ is chordal, we have $G_{i_0} = a - x_1 - x_2 - b$. Now we for any $1 \leq i \leq \ell$ and $i \neq i_0$ set the graph, $H_i := G_i / \{a, b\}$ if $a, b \in V(G_i)$ and $H_i := G_i$, otherwise. Therefore $H_1, H_2, \dots, H_{i_0-1}, H_{i_0+1}, \dots, H_\ell$ is a partition of edges of G and each H_i^c is chordal by Lemma 2.1. Hence we conclude that

$$\text{cochord}(G) \leq \text{cochord}(\mathcal{P}^4(G)) - 1.$$

Since $E(G_i)$ s cover the edges of $\mathcal{P}^4(G)$, without loss of generality, we can assume $x_2b \in E(G_1)$ and $x_1a \in E(G_2)$. If $x_1x_2 \notin E(G_1) \cup E(G_2)$, then there exists $3 \leq i \leq \ell$, such that $x_1x_2 \in E(G_i)$. Since G_i^c is a chordal graph, we conclude that $G_i \cong K_2$. Now set

$$H_1 := G_1 - x_2, H_2 := G_2 - x_1.$$

We conclude that the edge sets of subgraphs

$$H_1, H_2, G_3, \dots, G_{i-1}, G_{i+1}, \dots, G_\ell$$

is a partition of edges of G and their complements are chordal. Therefore

$$\text{cochord}(G) \leq \text{cochord}(\mathcal{P}^4(G)) - 1.$$

If $x_1x_2 \in E(G_1)$, then $N_{G_1}(b)$ is an independent set, because the complement of G_2 is a chordal graph. Suppose that $N_{G_1}(b) = \{x_2, w_1, w_2, \dots, w_t\}$ and

consider the graph

$$H := (\mathcal{P}^4(G) - \{bw_1, bw_2, \dots, bw_t\}) + \{aw_1, aw_2, \dots, aw_t\}.$$

Set

$$H_1 := a - x_1 - x_2 - b, H_2 := G_2 + \{aw_1, aw_2, \dots, aw_t\}$$

and $H_i := G_i$ for $3 \leq i \leq \ell$. It is not difficult to see that the edges of subgraphs H_1, H_2, \dots, H_ℓ is a partition of edges of H and H_i^c are chordal for $3 \leq i \leq \ell$. Hence the edges of subgraphs H_2, H_3, \dots, H_ℓ is a partition of edges of G and we conclude that

$$\text{cochord}(G) \leq \text{cochord}(H) \leq \text{cochord}(\mathcal{P}^4(G)) - 1.$$

□

3. Open problems

Problem 3.1. *Prove or disprove: Let G be a graph with $\text{cochord}(G) = \ell$. For any $x \in V(G)$, there exist a cochordal cover G_1, \dots, G_ℓ , such that there exists $1 \leq i \leq \ell$ which $xv \in E(G_i)$, for any $v \in N(x)$.*

Problem 3.2. *For any graph G ,*

$$\text{cochord}(\mathcal{P}^4(G)) \leq \text{cochord}(G) + 2.$$

If the statement of Problem 3.1 is true, then we can solve the Problem 3.2. Suppose that $\text{cochord}(G) = \ell$ and G_1, \dots, G_ℓ be edge disjoint subgraphs of G such that $E(G) = \bigcup_1^\ell E(G_i)$, and G_i^c are chordal graphs. Without loss of generality we can assume that $xv \in E(G_1)$ for any $v \in N(x)$. Hence $\mathcal{P}^4(G_1), G_2, \dots, G_\ell$, partitioned the edges of $\mathcal{P}^4(G)$. Now since G_1 is a chordal graph, we conclude that $\text{cochord}(\mathcal{P}^4(G)) \leq \text{cochord}(G) + 2$, by applying Lemma 2.2. Also if the statement of Problem 3.1 is true, then by applying Lemma 2.4, the following lemma is true.

Lemma 3.3. *Let G be a graph. Therefore $\text{cochord}(\mathcal{P}^4(G, x)) = \text{cochord}(G) + 1$, if one of the following holds:*

- i) One of the partition of $N(x)$ is empty;*
- ii) Induced subgraph $\langle N_G(x) \rangle$ is a clique.*

There are some families of graphs such as, complete graph, paths and cycle of length at least five, with the property $\text{cochord}(\mathcal{P}^4(G, x)) = \text{cochord}(G) + 1$. Hence it is natural to state the following problem.

Problem 3.4. *Characterize all graphs with property that for each vertex x and each partitions of $N_G(x)$, we have $\text{cochord}(\mathcal{P}^4(G, x)) = \text{cochord}(G) + 1$*

Problem 3.5. *Characterize all graphs G with $\text{cochord}(\mathcal{P}^4(G, x)) = 2$.*

Problem 3.6. *Characterize all graphs G for which $\text{cochord}(\mathcal{P}^4(G)) = 2$.*

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